

- 1 Introduction
- 2 C++ Basics
- 3 Data Structures
- 4 Application
- 5 Q & A

Self-Introduction



Information

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- **Year** Second-year Undergraduate
- **Major** Computer Science

Achievements

- **Silver Medal**, 50th ICPC Asia Regional Contest (Wuhan)
- **Silver Medal**, 50th ICPC Asia Regional Contest (Nanjing)
- **Bronze Medal**, 11th CCPC National Contest (Chongqing)

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Prec.	Operator	Assoc.
1	() [] -> . :: ++ -- (postfix)	L→R
2	! ~ ++ -- (prefix) - + * &	R→L
3	->* .*	L→R
4	* / %	L→R
5	+ -	L→R
6	« »	L→R
7	< <= > >=	L→R
8	== !=	L→R
9	&	L→R
10	^	L→R
11		L→R
12	&&	L→R
13		L→R
14	? :	R→L
15	= += -= *= /= %= &= ^= = «= »=	R→L
16	,	L→R

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Example

```
1 void f1 ()
2 {
3     int *p = new int[5];
4     int variable = 10;
5 }
```

```

1 0000000000001149 <_Z2f1v>:a
2      1149:  f3 0f 1e fa                                endbr64
3      114d:  55                                           push    %rbp
4      114e:  48 89 e5                                    mov     %rsp,%rbp
5      1151:  48 83 ec 10                                sub     $0x10,%rsp
6      1155:  bf 14 00 00 00                            mov     $0x14,%edi
7      115a:  e8 f1 fe ff ff                            call    1050 <_Znam@plt>
8      115f:  48 89 45 f8                                mov     %rax,-0x8(%rbp)
9      1163:  c7 45 f4 0a 00 00 00                      movl    $0xa,-0xc(%rbp)
10     116a:  90                                           nop
11     116b:  c9                                           leave
12     116c:  c3                                           ret

```

Example (Cont.)

```
1 void f1 ()  
2 {  
3     int *p = new int[5];  
4     delete[] p;  
5 }
```

```
1 00000000000001149 <_Z2f1v>:  
2 1169: f3 0f 1e fa          endbr64  
3 116d: 55                    push    %rbp  
4 116e: 48 89 e5              mov     %rsp,%rbp  
5 1171: 48 83 ec 10          sub     $0x10,%rsp  
6 1175: bf 14 00 00 00       mov     $0x14,%edi  
7 117a: e8 e1 fe ff ff       call   1060 <_Znam@plt>  
8 117f: 48 89 45 f8          mov     %rax,-0x8(%rbp)  
9 1183: 48 83 7d f8 00       cmpq   $0x0,-0x8(%rbp)  
10 1188: 74 0c               je      1196 <_Z2f1v+0x2d>  
11 118a: 48 8b 45 f8          mov     -0x8(%rbp),%rax  
12 118e: 48 89 c7            mov     %rax,%rdi  
13 1191: e8 da fe ff ff       call   1070 <_Zdapv@plt>  
14 1196: 90                  nop  
15 1197: c9                  leave  
16 1198: c3                  ret
```

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```
1 | #include <stdio>
2 |
3 | void calc (int a,int b,int **f1,int **f2)
4 | {
5 |     **f1 = a + b;**f2 = a - b;
6 |     int *temp = *f1;*f1 = *f2;*f2 = temp;
7 | }
8 |
9 | int main ()
10 | {
11 |     int a = 5,b = 3,x,y;
12 |     int *px = &x,*py = &y;
13 |     calc (a,b,&px,&py);
14 |     printf ("%d□%d\n",x,y);
15 |     return 0;
16 | }
```

8 2

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Array (Cont.)

```
1 #include <stdio>
2 /* declaration here */
3 int main ()
4 {
5     char arr[2][3] = {{ 'a', 'b', 'c' }, { 'e', 'f', 'g' }};
6     calc (arr);
7     return 0;
8 }
```

Which of the following function declarations are correct?

- A) void calc (char (*arr)[3]);
- B) void calc (char arr[2][3]);
- C) void calc (char **arr);
- D) void calc (char arr[][3]);

Array (Cont.)

```
1 #include <stdio>
2 /* declaration here */
3 int main ()
4 {
5     char arr[2][3] = {{ 'a', 'b', 'c' }, { 'e', 'f', 'g' }};
6     calc (arr);
7     return 0;
8 }
```

Which of the following function declarations are correct?

- A) void calc (char (*arr)[3]);
- B) void calc (char arr[2][3]);
- C) void calc (char **arr);
- D) void calc (char arr[][3]);

Answer

1 ABD


```
1 #include <stdio>
2 using namespace std;
3 int main ()
4 {
5     int a[] = {3,1,4,1,5,9};
6     int *p = a + 1,*q = a + 4;
7     int x = ++*p,y = *q--;
8     int z = ++*p,w = (*q)--;
9     int t = (*p++) + (*--q);
10    int u = *p,v = *q;
11    for (int i = 0;i < 6;++i) printf ("%d",a[i]);
12    printf ("\n%d%d%d%d%d%d%d\n",x,y,z,w,t,u,v);
13    return 0;
14 }
```


Vector

```
1  int main ()
2  {
3
4      vector <int> lst = {2,0,2,5,1,2,1,4};
5      int tot = 0, flag = 0, n = lst.size ();
6      for (auto it = lst.begin (); it != lst.end (); ++it)
7      {
8          if (flag) tot += lst.front ();
9          else tot += *it;
10         flag = 1 - flag;
11     }
12     printf ("%d,%d\n", tot, tot / n);
13     return 0;
14 }
```

Vector

```
1  int main ()
2  {
3      vector<int> lst = {2,0,2,5,1,2,1,4};
4      int tot = 0, flag = 0, n = lst.size ();
5      for (auto it = lst.begin (); it != lst.end (); ++it)
6      {
7          if (flag) tot += lst.front ();
8          else tot += *it;
9          flag = 1 - flag;
10     }
11     printf ("%d,%d\n", tot, tot / n);
12     return 0;
13 }
```

Answer

1 14,1

Vector (Cont.)

resize (),reserve ()

```
1 | vector <int> lst = {2,0,2,5};  
2 | lst.reserve (8);  
3 | printf ("%ld,%ld,%d\n",lst.size (),lst.capacity (),lst.back ());  
4 | lst.reserve (6);  
5 | printf ("%ld,%ld,%d\n",lst.size (),lst.capacity (),lst.back ());  
6 | lst.resize (10,3);  
7 | printf ("%ld,%ld,%d\n",lst.size (),lst.capacity (),lst.back ());  
8 | lst.resize (3,5);  
9 | printf ("%ld,%ld,%d\n",lst.size (),lst.capacity (),lst.back ());
```

Vector (Cont.)

resize (),reserve ()

```
1 | vector <int> lst = {2,0,2,5};  
2 | lst.reserve (8);  
3 | printf ("%ld,%ld,%d\n",lst.size (),lst.capacity (),lst.back ());  
4 | lst.reserve (6);  
5 | printf ("%ld,%ld,%d\n",lst.size (),lst.capacity (),lst.back ());  
6 | lst.resize (10,3);  
7 | printf ("%ld,%ld,%d\n",lst.size (),lst.capacity (),lst.back ());  
8 | lst.resize (3,5);  
9 | printf ("%ld,%ld,%d\n",lst.size (),lst.capacity (),lst.back ());
```

Answer

```
1 | 4,8,5  
2 | 4,8,5  
3 | 10,10,3  
4 | 3,10,2
```



```
1 int main ()
2 {
3     CSC arr;
4     for (int i = 1; i <= 5; ++i) arr.add (i);
5     for (int i = 0; i <= 5; ++i)
6     {
7         printf ("%d", arr.del ());
8         if (i != 5) printf (",");
9     }
10    printf ("\n");
11    return 0;
12 }
```



```
1 int main ()
2 {
3     CSC arr;
4     for (int i = 1; i <= 5; ++i) arr.add (i);
5     for (int i = 0; i <= 5; ++i)
6     {
7         printf ("%d", arr.del ());
8         if (i != 5) printf (",");
9     }
10    printf ("\n");
11    return 0;
12 }
```

Answer

1, 2, 3, 4, 5, -1

Stack/Queue (Cont.)

Sliding Window

```
1 void calc ()
2 {
3     int q[1005] = {0}, num[] = {1,2,7,3,8,5,2,9};
4     int k = 3, n = 8, head = 0, tail = -1;
5     for (int i = 0; i < n; ++i)
6     {
7         while (head <= tail && q[head] <= i - k) ++head;
8         while (head <= tail && num[q[tail]] <= num[i]) --tail;
9         q[++tail] = i;
10        if (i >= k - 1) printf ("%d", num[q[head]]);
11    }
12    printf ("\n");
13 }
```

Stack/Queue (Cont.)

Sliding Window

```
1 void calc ()
2 {
3     int q[1005] = {0}, num[] = {1,2,7,3,8,5,2,9};
4     int k = 3, n = 8, head = 0, tail = -1;
5     for (int i = 0; i < n; ++i)
6     {
7         while (head <= tail && q[head] <= i - k) ++head;
8         while (head <= tail && num[q[tail]] <= num[i]) --tail;
9         q[++tail] = i;
10        if (i >= k - 1) printf ("%d", num[q[head]]);
11    }
12    printf ("\n");
13 }
```

Answer (Hint: It also can be implemented with `std::deque`.)

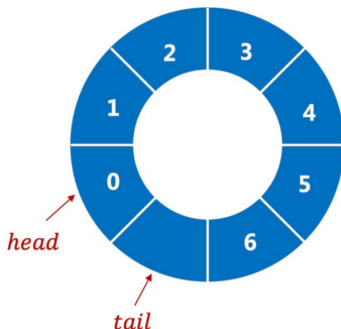
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Stack/Queue (Cont.)

Circular Queue

Empty $\text{head} == \text{tail};$

Full $(\text{tail} + 1) \% \text{capacity} == \text{front};$



Link list

Version 1

```
1  LinkNode* func (LinkNode* head)
2  {
3      LinkNode* pre = nullptr;
4      LinkNode* cur = head;
5      LinkNode* nxt = nullptr;
6      while (cur != nullptr)
7      {
8          nxt = cur -> next;
9          cur -> next = pre;
10         pre = cur; cur = nxt;
11     }
12     return pre;
13 }
```

Link list (Cont.)

Version 2

```
1 LinkNode* func (LinkNode* head)
2 {
3     if (head == nullptr || head -> next == nullptr) return head;
4     LinkNode* new_head = func (head -> next);
5     head -> next -> next = head;
6     head -> next = nullptr;
7     return new_head;
8 }
```

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Sequence Containers

Associative Containers

Tree

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Map

Hashmap

```
1 #define TABLE_SIZE 7
2 struct HashTable {int table[TABLE_SIZE];};
3 int hashFunction (int key) {return key % TABLE_SIZE;}
4 void insert(int table[],int key)
5 {
6     int id = hashFunction (key);
7     while (table[id] != -1) id = (id + 1) % TABLE_SIZE;
8     table[id] = key;
9 }
```


Map (Cont.)

1. The following keys are inserted in order: {7, 14, 21, 28, 35, 42}. After inserting all keys, what is the average successful search length (average number of comparisons to find an existing key)?
 - A) 2.0
 - B) 2.5
 - C) 3.0
 - D) 3.5
2. In the worst-case scenario for linear probing, inserting n elements into a table of size n requires?
 - A) $O(n)$ time per insertion
 - B) $O(\log n)$ time per insertion
 - C) $O(n \log n)$ total time for all insertions
 - D) $O(n^2)$ total time for all insertions

Answer

1

D ; AD

Set

Bitset

```
1  const int RANGE_SIZE = 128;
2  const int BITS_PER_BYTE = 8;
3  const int BITS_PER_LONG = BITS_PER_BYTE * sizeof(long);
4  const int CVEC_WORDS = (RANGE_SIZE + BITS_PER_LONG - 1) / BITS_PER_LONG;
5  struct BitSet
6  {
7      unsigned long words[CVEC_WORDS];
8      BitSet ()
9      {
10         for (int i = 0; i < CVEC_WORDS; ++i) words[i] = 0;
11     }
12 };
13 unsigned long createMask (int k)
14 {
15     unsigned long x = 1L;
16     return x << k % BITS_PER_LONG;
17 }
```

Set (Cont.)

```
1 // check if the k-th characters in ASCII table in the Set
2 bool inSet(BitSet &cv, int k)
3 {
4     if (k < 0 || k >= RANGE_SIZE) return;
5     return cv.words[k / BITS_PER_LONG] & createMask (k);
6 }
7 // set the k-th character in ASCII table to the set
8 void setBit(BitSet &cv, int k)
9 {
10    if (k < 0 || k >= RANGE_SIZE) return;
11    cv.words[k / BITS_PER_LONG] |= createMask (k);
12 }
13 // remove the k-th characters in ASCII Table from the set
14 void remove(BitSet &cv, int k)
15 {
16    if (k < 0 || k >= RANGE_SIZE) return;
17    cv.words[k / BITS_PER_LONG] &= ~createMask (k);
18 }
```

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Binary Tree

Three Tree Traversal Methods

- **Preorder Traversal** Root → Left → Right
- **Inorder Traversal** Left → Root → Right
- **Postorder Traversal** Left → Right → Root

Question

Preorder traversal 1, 2, 4, 8, 5, 9, 3, 6, 10, 7

Inorder traversal 8, 4, 2, 9, 5, 1, 6, 10, 3, 7

Based on the given preorder and inorder traversals, what is the postorder traversal?

Binary Tree

Three Tree Traversal Methods

- **Preorder Traversal** Root → Left → Right
- **Inorder Traversal** Left → Root → Right
- **Postorder Traversal** Left → Right → Root

Question

Preorder traversal 1, 2, 4, 8, 5, 9, 3, 6, 10, 7

Inorder traversal 8, 4, 2, 9, 5, 1, 6, 10, 3, 7

Based on the given preorder and inorder traversals, what is the postorder traversal?

Answer

8, 4, 9, 5, 2, 10, 6, 7, 3, 1

```

1 | int find (int x)
2 | {
3 |     for (int i = 0; i < n; ++i)
4 |         if (in[i] == x) return i;
5 |     return -1;
6 | }
7 | void getpost (int sl, int sr, int fl, int fr)
8 | {
9 |     if (sl > sr || fl > fr) return;
10 |    int k = find (__1__);
11 |    getpost (__2__); getpost (__3__);
12 |    printf ("%d\n", pre[sl]);
13 | }
14 | void solve () {getpost (0, n - 1, 0, n - 1);}

```


Binary Tree (Cont.)

```
1 int find (int x)
2 {
3     for (int i = 0; i < n; ++i)
4         if (in[i] == x) return i;
5     return -1;
6 }
7 void getpost (int sl, int sr, int fl, int fr)
8 {
9     if (sl > sr || fl > fr) return;
10    int k = find (__1__);
11    getpost (__2__); getpost (__3__);
12    printf ("%d\n", pre[sl]);
13 }
14 void solve () {getpost (0, n - 1, 0, n - 1);}
```

Answer

```
1 pre[sl]
2 sl + 1, sl + k - fl, fl, k - 1
3 sl + k - fl + 1, sr, k + 1, fr
```

BST

```
1 struct TreeNode
2 {
3     int key,sz,cnt;
4     TreeNode *left,*right;
5     TreeNode (int val) :
6         key (val), sz (1), cnt (1), left (nullptr), right (nullptr) {}
7 };
```

```
1 bool search (TreeNode* u,int val)
2 {
3     if (u == nullptr) return false;
4     if (u -> key == val) return true;
5     else if (val < u -> key) return search (rt -> left,val);
6     else return search (rt -> right,val);
7 }
```

BST (Cont.)

```
1 int get_sz (TreeNode* u) {return u == nullptr ? 0 : u -> size;}
2
3
4
5
6
7
8
9
10 }
```

```
1
2
3
4
5
TreeNode* findMinNode (TreeNode* u)
{
    while (u -> left != nullptr) u = u -> left;
    return u;
}
```

BST (Cont.)

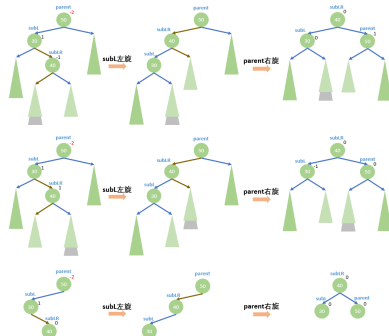
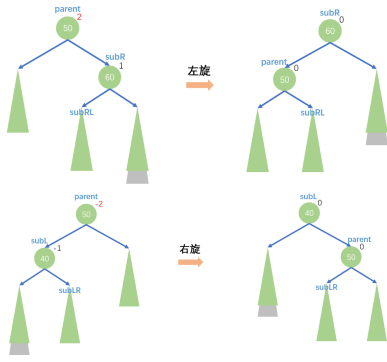
```
1 | TreeNode* remove (TreeNode* u,int val)
2 | {
3 |
4 |     if (u == nullptr) return u;
5 |     if (val < u -> key) u -> left = remove (u -> left,val);
6 |     else if (val > u -> key) u -> right = remove (u -> right,val);
7 |     else
8 |     {
9 |         if (u -> cnt > 1) {--(u -> cnt);return;}
10 |        if (u -> left != nullptr && u -> right != nullptr)
11 |        {
12 |            TreeNode* nxt = findMinNode (u -> right);
13 |            u -> key = nxt -> key;u -> cnt = nxt -> cnt;nxt -> cnt = 1;
14 |            u -> right = remove (u -> right,nxt -> key);
15 |        }
16 |        else
17 |        {
18 |            TreeNode *tmp = u;
19 |            u = (u -> left != nullptr) ? u -> left : u -> right;
20 |            delete tmp;return u;
21 |        }
22 |    }
23 |    u -> sz = u -> cnt + get_sz (u -> left) + get_sz (u -> right);
24 |    return u;
25 | }
```

BST (Cont.)

```
1 int queryRank (TreeNode* u, int val)
2 {
3     if (u == nullptr) return;
4     if (v == u -> key) return get_sz (u -> left) + 1;
5     else if (v < u -> key) return queryRank (u -> left, val);
6     return queryRank (u -> right, val) + get_sz (u -> left) + u -> cnt;
7 }
```

```
1 int querykth (TreeNode* u, int k)
2 {
3     if (u == nullptr) return -1;
4     if (u -> left)
5     {
6         if (u -> left -> sz >= k) return querykth (u -> left, k);
7         if (u -> left -> sz + u -> cnt >= k) return u -> key;
8     }
9     else if (k <= u -> cnt) return u -> cnt;
10    return querykth (u -> right, k - get_sz (u -> left) - u -> cnt);
11 }
```

AVL Tree



AVL Tree (Cont.)

```
1 | avl_node* R_rotate (avl_node* u)
2 | {
3 |     avl_node* x = u -> L; avl_node* y = x -> R;
4 |     x -> R = u; u -> L = y;
5 |     x -> fat = u -> fat;
6 |     u -> fat = x;
7 |     if (y != nullptr) y -> fat = u;
8 |     u -> upd (); x -> upd ();
9 |     return x;
10 | }
11 | avl_node* L_rotate (avl_node* u)
12 | {
13 |     avl_node* x = u -> R; avl_node* y = x -> L;
14 |     x -> L = u; u -> R = y;
15 |     x -> fat = u -> fat;
16 |     u -> fat = x;
17 |     if (y != nullptr) y -> fat = u;
18 |     u -> upd (); x -> upd ();
19 |     return x;
20 | }
```

AVL Tree (Cont.)

```

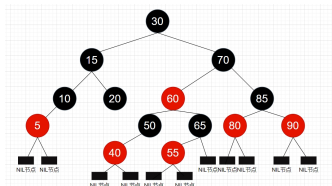
1 | avl_node* LR_rotate (avl_node* u) {u -> L = L_rotate (u -> L);return R_rotate (u);}
2 | avl_node* RL_rotate (avl_node* u) {u -> R = R_rotate (u -> R);return L_rotate (u);}
3 |
4 | int height (avl_node* u) const {return u ? u -> dep : -1;}
5 | int calc (avl_node* u) const {return !u ? 0 : height (u -> L) - height (u -> R);}
6 | avl_node* balance (avl_node* u)
7 | {
8 |     if (u == nullptr) return nullptr;
9 |     int fac = calc (u);
10 |    if (fac > 1)
11 |    {
12 |        if (calc (u -> L) >= 0) return R_rotate (u);
13 |        u -> L = L_rotate (u -> L);
14 |        return R_rotate (u);
15 |    }
16 |    else if (fac < -1)
17 |    {
18 |        if (calc (u -> R) <= 0) return L_rotate (u);
19 |        u -> R = R_rotate (u -> R);
20 |        return L_rotate (u);
21 |    }
22 |    return u;
  }

```


Red-Black Tree

Key Characteristics of RB Trees

- Binary Search Tree Property
- Color Rules
 - Every node is either **red** or **black**
 - Root is always **black**
 - All leaves (NIL nodes) are **black**
- Red Node Constraint
 - No two consecutive red nodes
- Black Height Property
 - Every path from root to leaf has the same number of black nodes
 - Ensures tree remains balanced



Red-Black Tree (Cont.)

Theorem

For a red-black tree with n nodes, its height h satisfies
 $h \leq 2 \log_2(n + 1)$.

Proof

Let bh be the black height of a RB tree.
Then we have $n \geq 2^{bh} - 1$ and $h \leq 2bh$.
So we can conclude $h \leq 2bh \leq \log_2(n + 1)$.

Heap

Heap Construction

$$\begin{aligned}T(n) &= 1 \times (H-1) + 2 \times (H-2) + \dots + 2^{H-1} \times 0 \\&= 2^H - H - 1 = O(n)\end{aligned}$$

```
1 void down (int *a,int n,int i)
2 {
3     int pos = i,l = 2 * i + 1,r = 2 * i + 2;
4     if (l < n && a[l] > a[pos]) pos = l;
5     if (r < n && a[r] > a[pos]) pos = r;
6     if (pos != i) swap (a[i],a[pos]),down (a,n,pos);
7 }
8 void build () //Default: max-heap
9 {
10     for (int i = a.size () / 2 - 1;i >= 0;--i) down (a,a.size (),i);
11 }
```

Heap (Cont.)

Insert

```
1 void up (int *a,int id)
2 {
3     int fa = (id - 1) / 2;
4     while (id > 0 && a[id] > a[fa])
5         {swap (a[id],a[fa]);id = fa;fa = (id - 1) / 2;}
6 }
7 void insert (int val)
8 {
9     if (n == _MAX_SIZE) return;
10    a[n] = val;++n;up (a,n);
11 }
```

Delete

```
1 void del (int *a,int n)
2 {
3     if (n <= 0) return;
4     swap (a[0],a[n]);--n;down (a,n,0);
5 }
```

Heap (Cont.)

Q1. Given the keyword sequence 5, 8, 12, 19, 28, 20, 15, 22 is a **min-heap**, after inserting keyword 3 and performing heap adjustment, the resulting min-heap is:

- A) 3, 5, 12, 8, 28, 20, 15, 22, 19
- B) 3, 5, 12, 19, 20, 15, 22, 8, 28
- C) 3, 8, 12, 5, 20, 15, 22, 28, 19
- D) 3, 12, 5, 8, 28, 20, 15, 22, 19

Heap (Cont.)

Q1. Given the keyword sequence 5, 8, 12, 19, 28, 20, 15, 22 is a **min-heap**, after inserting keyword 3 and performing heap adjustment, the resulting min-heap is:

- A) 3, 5, 12, 8, 28, 20, 15, 22, 19
- B) 3, 5, 12, 19, 20, 15, 22, 8, 28
- C) 3, 8, 12, 5, 20, 15, 22, 28, 19
- D) 3, 12, 5, 8, 28, 20, 15, 22, 19

Answer

1

A

Heap (Cont.)

Q2. Which of the following statements about a **max-heap** (containing at least 2 elements) is/are correct?

- I. A heap can be viewed as a **complete binary tree**.
 - II. A heap can be stored using **sequential (array-based) storage**.
 - III. A heap can be viewed as a **binary search tree**.
 - IV. The **second largest element** in the heap must be in the level directly below the root.
- A) I and II only
B) II and III only
C) I, II, and IV only
D) I, III, and IV only

Heap (Cont.)

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Answer

1

c

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Complexity

Complexity	Name	Example
$O(1)$	Constant	Array access
$O(\log n)$	Logarithmic	Binary search
$O(n)$	Linear	Linear search
$O(n \log n)$	Log-linear	Merge sort
$O(n^2)$	Quadratic	Bubble sort
$O(n^3)$	Cubic	Matrix multiplication
$O(2^n)$	Exponential	Subset enumeration

Growth Rates:

$$O(1) \leq O(\log n) \leq O(n) \leq O(n \log n) \leq O(n^2) \leq O(n^3) \leq O(2^n)$$

Complexity (Cont.)

Time Complexity Analysis

```
1 for (int i = 1; i <= n; ++i)
2     for (int j = 1; j <= n / 2; ++j) ++cnt;
3 for (int i = 1; i <= n; i *= 2)
4     for (int j = 1; j <= n; ++j) ++cnt;
5 for (int i = 1; i <= n; i *= 2)
6     for (int j = 1; j <= i; ++j) ++cnt;
7 for (int i = 1; i <= n; ++i)
8     for (int j = 1; j <= n / i; ++j) ++cnt;
```

Complexity (Cont.)

Time Complexity Analysis

```
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5 for (int i = 1; i <= n; i *= 2)
6     for (int j = 1; j <= i; ++j) ++cnt;
7 for (int i = 1; i <= n; ++i)
8     for (int j = 1; j <= n / i; ++j) ++cnt;
```

Answer

$O(n^2)$, $O(n \log n)$, $O(n)$, $O(n \log n)$

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Sort

```
1 void Bubble_Sort () {  
2     for (int i = 0; i < n;++i)  
3         for (int j = 0;j < n - i - 1;++j)  
4             if (a[j] > a[j + 1]) swap (a[j],a[j + 1]);  
5 }  
6 void Selection_Sort () {  
7     for (int i = 0;i < n;++i)  
8     {  
9         int minIndex = i;  
10        for (int j = i + 1;j < n;++j)  
11            if (a[j] < a[minIndex]) minIndex = j;  
12        swap (a[minIndex],a[i]);  
13    }  
14 }  
15 void Insertion_Sort () {  
16     for (int i = 1;i < n;++i)  
17     {  
18         int key = a[i],j = i - 1;  
19         while (j >= 0 && a[j] > key) a[j + 1] = a[j],--j;  
20         a[j + 1] = key;  
21     }  
22 }
```

Sort (Cont.)

```
1  template <typename T>
2  void Merge_Sort (T a[],int l,int r)
3  {
4      if (l >= r) return;
5      int mid = (l + r) / 2;
6      Merge_Sort (a,l,mid);Merge_Sort (a,mid + 1,r);
7      vector <T> lst;
8      int posl = l,posr = mid + 1;
9      while (posl <= mid && posr <= r)
10     {
11         if (a[posl] <= a[posr]) lst.push_back (a[posl++]);
12         else lst.push_back (a[posr++]);
13     }
14     while (posl <= mid) lst.push_back (a[posl++]);
15     while (posr <= r) lst.push_back (a[posr++]);
16     for (int i = l;i <= r;++i) a[i] = lst[i - l];
17 }
18 Merge_Sort (a,0,n - 1);
```


Sort (Cont.)

```
1  template <typename T>
2  int partition (T a[],int l,int r)
3  {
4      int pivot = a[l];
5      while (l < r)
6      {
7          while (l < r && pivot <= a[r]) --r;
8          a[l] = a[r];
9          while (l < r && a[l] <= pivot) ++l;
10         a[r] = a[l];
11     }
12     a[l] = pivot;
13     return l;
14 }
15 template <typename T>
16 void Quick_Sort (T a[],int l,int r)
17 {
18     if (l >= r) return;
19     int pivot = partition (a,l,r);
20     Quick_Sort (a,l,pivot - 1);Quick_Sort (a,pivot + 1,r);
21 }
22 Quick_Sort (a,0,n - 1);
```

Sort (Cont.)

The Complexity of Different Sort Algorithms

Algorithm	Best	Average	Worst
Selection Sort	$O(n^2)$	$O(n^2)$	$O(n^2)$
Insertion Sort	$O(n)$	$O(n^2)$	$O(n^2)$
Bubble Sort	$O(n)$	$O(n^2)$	$O(n^2)$
Merge Sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
Quick Sort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$
Heap Sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
BST Tree Sort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$

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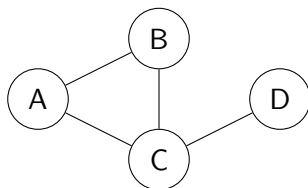
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Graph



Adjacency List : $O(|V| + |E|)$.

A: → B → C → ×

B: → A → C → ×

C: → A → B → D → ×

D: → C → ×

Adjacency Matrix : $O(|V|^2)$.

$$\begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

DFS/BFS

Visualization

Algorithm 1 DFS

```
procedure SEARCH( $u$ )  
    mark  $u$  as vis  
    for each neighbor  $v$  of  $u$  do  
        if  $v$  is not vis then  
            SEARCH( $v$ )  
        end if  
    end for  
end procedure
```

Algorithm 2 BFS

```
procedure SEARCH( $source$ )  
     $Q \leftarrow$  new Queue()  
     $Q.enqueue(source)$   
    while  $Q$  not empty do  
         $u \leftarrow Q.dequeue()$   
        for each neighbor  $v$  of  $u$  do  
            if  $v$  is not vis then  
                mark  $v$  as vis  
                 $Q.enqueue(v)$   
            end if  
        end for  
    end while  
end procedure
```

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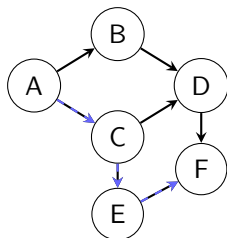
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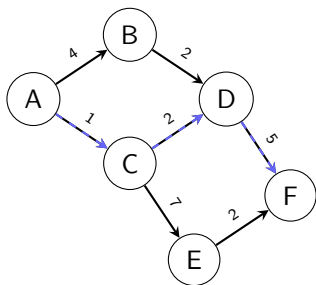
Unweighted Shortest-Path



Algorithm 3 0-1 BFS

```
procedure CALC(source, graph)  
     $Q \leftarrow \text{new Queue}()$   
     $\text{dist} \leftarrow \text{new Dictionary}()$   
     $Q.\text{enqueue}(\text{source})$   
     $\text{dist}[\text{source}] \leftarrow 0$   
    while  $Q$  is not empty do  
         $u \leftarrow Q.\text{dequeue}()$   
        for each neighbor  $v$  of  $u$  in graph do  
            if  $v$  is not in dist then  
                 $\text{dist}[v] \leftarrow \text{dist}[u] + 1$   
                 $Q.\text{enqueue}(v)$   
            end if  
        end for  
    end while  
    return dist  
end procedure
```

Weighted Shortest-Path



Algorithm 4 Dijkstra

```

procedure CALC(source, graph)
  dist  $\leftarrow$  array of size  $|V|$ , initialized to  $\infty$ 
  vis  $\leftarrow$  boolean array of size  $|V|$ , initialized to false
  dist[source]  $\leftarrow$  0
  for  $i \leftarrow 1$  to  $|V|$  do
    u  $\leftarrow$  vertex with minimum dist not vis
    vis[u]  $\leftarrow$  true
    for each neighbor v of u with weight w do
      if not vis[v] and dist[u] + w < dist[v] then
        dist[v]  $\leftarrow$  dist[u] + w
      end if
    end for
  end for
  return dist
end procedure
  
```


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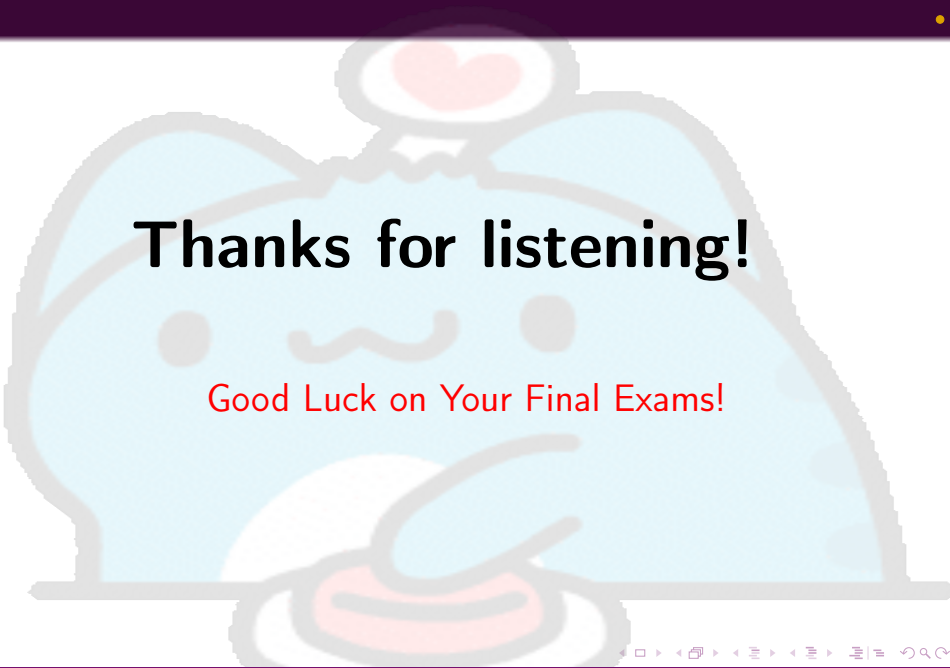
Algorithm 5 Topological Sort

```
procedure SEARCH( $G(V, E)$ )  
   $\text{indeg}[|V|] \leftarrow 0, q \leftarrow \emptyset, \text{ord} \leftarrow \emptyset$   
  for  $(u, v) \in E$  do  $\text{indeg}[v] \leftarrow \text{indeg}[v] + 1$   
  end for  
  for  $v \in V$  do  
    if  $\text{indeg}[v] = 0$  then  $q.\text{push}(v)$   
    end if  
  end for  
  while  $q \neq \emptyset$  do  
     $u \leftarrow q.\text{pop}(), \text{ord}.\text{push}(u)$   
    for  $v \in \text{adj}[u]$  do  
       $\text{indeg}[v] \leftarrow \text{indeg}[v] - 1$   
      if  $\text{indeg}[v] = 0$  then  $q.\text{push}(v)$   
      end if  
    end for  
  end while  
  if  $|\text{ord}| \neq |V|$  then return  $\emptyset$   
  end if  
  return  $\text{ord}$   
end procedure
```

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Q & A

Q & A



Thanks for listening!

Good Luck on Your Final Exams!